

Two New Unimodal Descent Polynomials

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Outline

Introduction and Main Results

Hierarchy and Implication

Main Results

Separable permutation, Schröder word and di-sk tree

Direct sum and skew sum

Sweeping-algorithm

The di-sk tree

$S_n(t)$ is γ -nonnegative

Steps in the proof

The bijection

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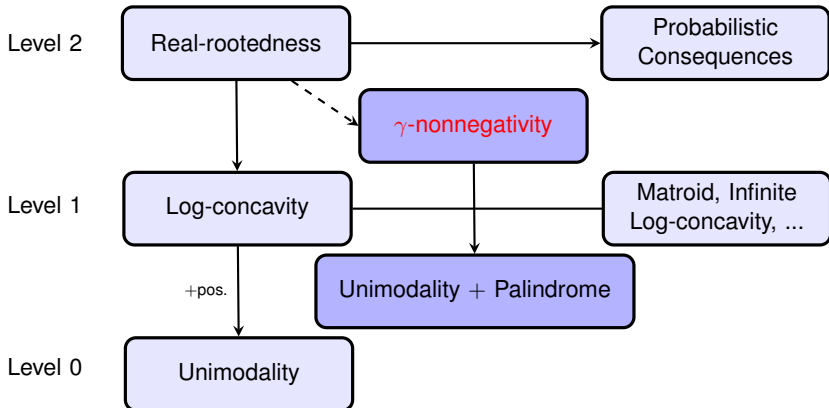
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Example

The n th row in Pascal's triangle $\{\binom{n}{k}\}_{k=0}^n$, i.e. $(1+x)^n$ satisfies all five properties above.

Hierarchy and More



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- Classical example includes Foata-Strehl’s “valley-hopping” for Eulerian polynomial.

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- This leads us naturally to the operations \oplus and \ominus , and I am about to tell the rest of the story...

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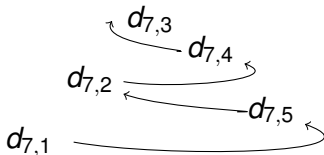
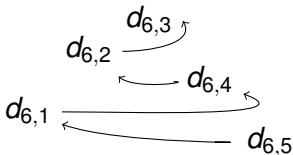
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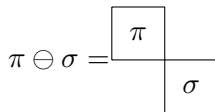
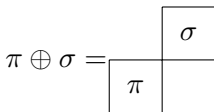
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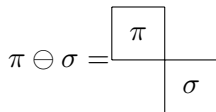
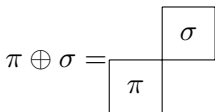


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- ii) We repeat this process until no new blocks can be formed.
- iii) The process will end with a single block if and only if it avoids patterns (3142) and (2413). We call such π a *separable permutation*, otherwise non-separable. For separable π , we replace all the numbers with 1 (since their order has been coded by \oplus , \ominus and parenthesis) and call the final expression, denoted by $\text{sw}(\pi)$, a *Schröder word*.

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sweep 5: $(9 \ominus 8)(4 \ominus (1 \oplus (3 \ominus 2)))7(5 \oplus 6)$

sweep 6: $(9 \ominus 8)(4 \ominus (1 \oplus (3 \ominus 2)))(7 \ominus (5 \oplus 6))$

sweep 7: $(9 \ominus 8)((4 \ominus (1 \oplus (3 \ominus 2))) \oplus (7 \ominus (5 \oplus 6)))$

sweep 8: $((9 \ominus 8) \ominus ((4 \ominus (1 \oplus (3 \ominus 2))) \oplus (7 \ominus (5 \oplus 6))))$

final: $((1 \ominus 1) \ominus ((1 \ominus (1 \oplus (1 \ominus 1))) \oplus (1 \ominus (1 \oplus 1)))) = \text{sw}(\pi)$

Right Chain Condition

Proposition (Shapiro? West? Folkloric?)

For $n \geq 1$, the sweeping-algorithm $\pi \mapsto \text{sw}(\pi)$ is a bijection between $\mathfrak{S}_n(2413, 3142)$ and \mathcal{W}_n .

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Remark

Associativity gives us $(\pi \oplus \sigma) \oplus \tau = \pi \oplus (\sigma \oplus \tau)$ and $(\pi \ominus \sigma) \ominus \tau = \pi \ominus (\sigma \ominus \tau)$. Since we always sweep from left to right in our decomposition, we will always first get block $(B_1 \oplus B_2) \oplus B_3$ before we can form $B_1 \oplus (B_2 \oplus B_3)$, similarly for operator \ominus .

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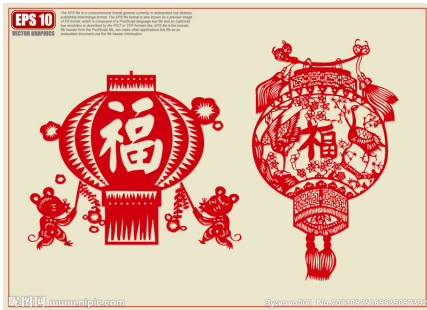
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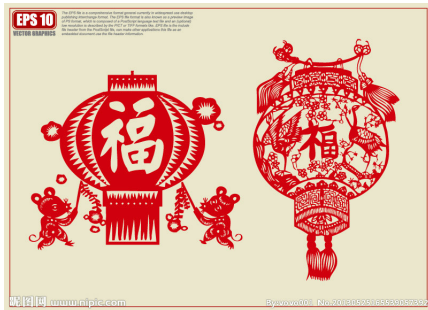




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Word to Tree

- Associate each pair of parenthesis in a Schröder word with the unique operator (\oplus or \ominus) it is “parenting”.

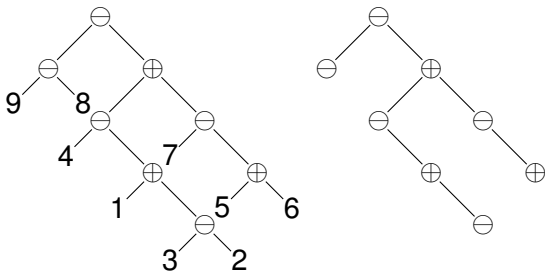
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Definition

Given a binary tree T , its *right chain* is any maximal chain composed of right children except the first node, which is either the root or a left child. The *length* of a right chain is the number of nodes on this chain. And the number of its right chains is denoted as $r(T)$. Similarly, the number of right chains with even (resp. odd) length is denoted as $r_e(T)$ (resp. $r_o(T)$).

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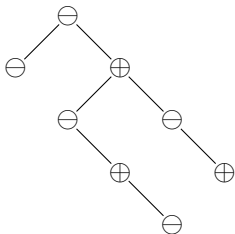
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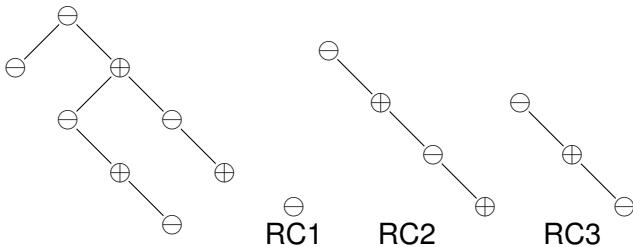
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We call a binary tree *di-sk tree* if it is labelled with \oplus and \ominus and the labelling satisfies the *right chain condition*. The set of all di-sk trees with $n - 1$ nodes is denoted as $\mathfrak{D}\mathfrak{T}_n$.

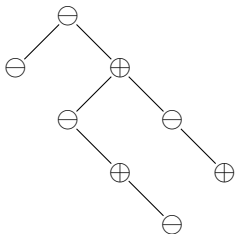
Right Chain, Lock and Hang



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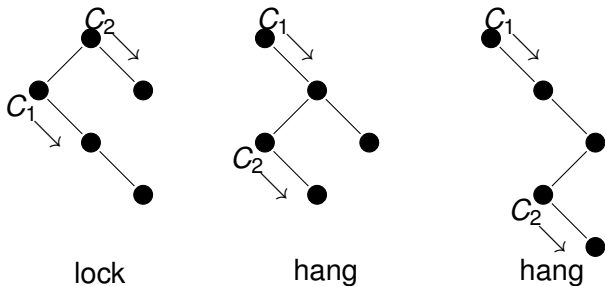
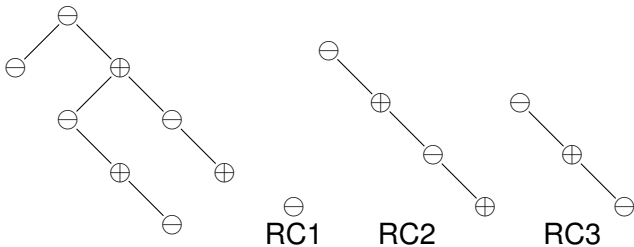
○
RC1

RC2

RC3

order?

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Back to see the example.

Outline

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Hierarchy and Implication

Main Results

Separable permutation, Schröder word and di-sk tree

Direct sum and skew sum

Sweeping-algorithm

The di-sk tree

$S_n(t)$ is γ -nonnegative

Steps in the proof

The bijection

Conclusion

Sketch of proof

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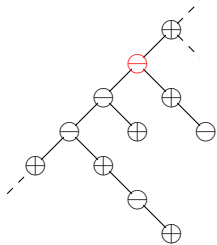
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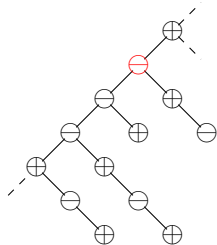
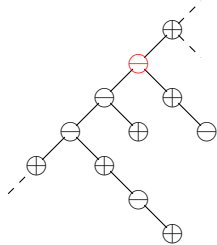
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- **Idea:** Cut and paste at “appropriate” places!

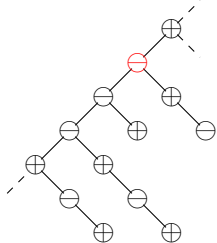
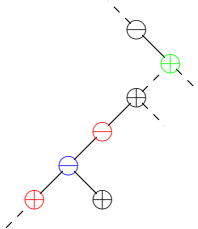
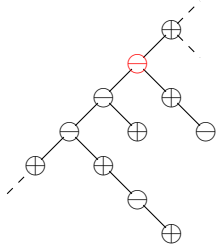
$$\mathcal{DT}^2 \setminus \mathcal{DT}^1 \xrightarrow{\psi} \mathcal{DT}^1 \setminus \mathcal{DT}^2$$



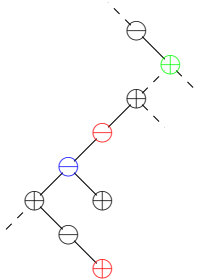
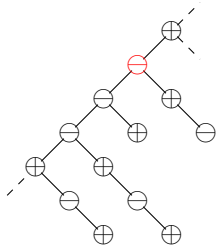
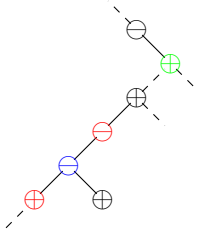
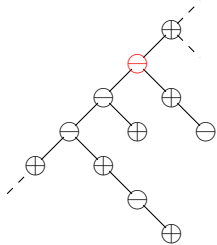
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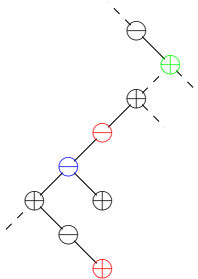
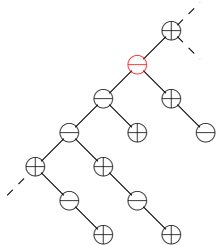
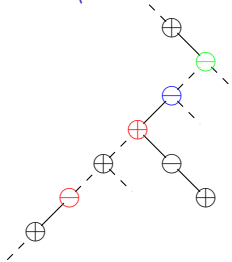
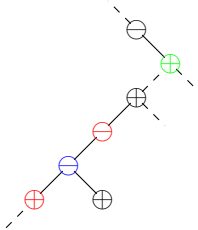
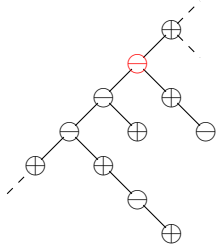
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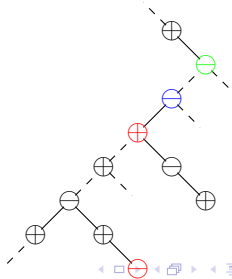
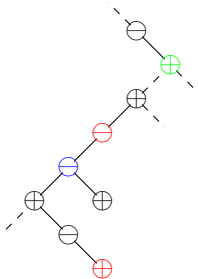
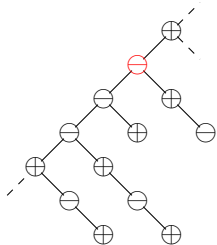
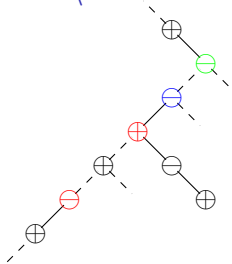
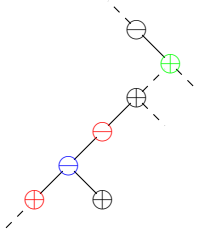
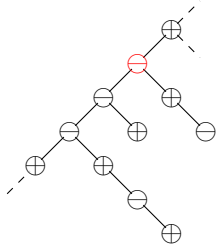
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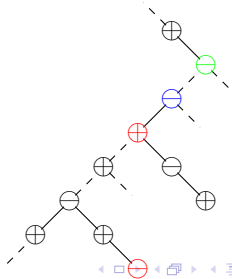
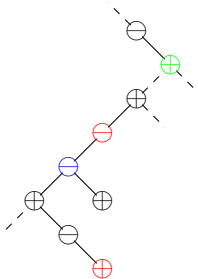
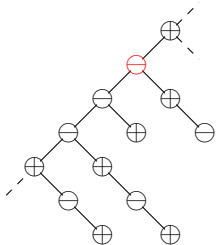
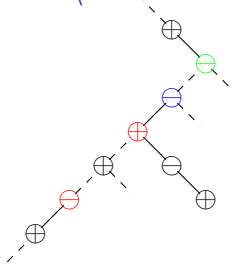
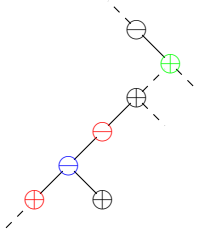
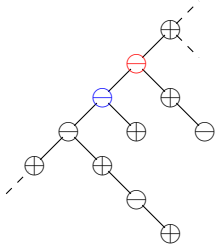
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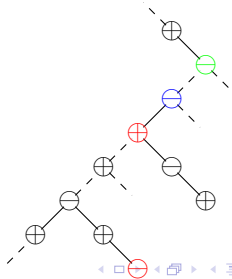
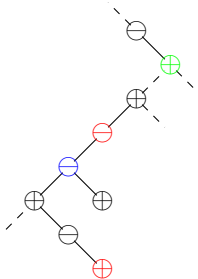
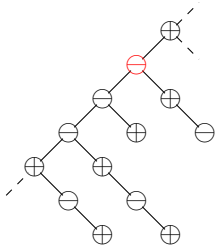
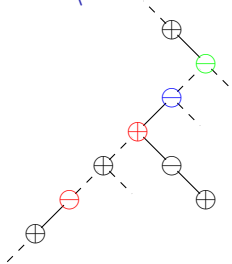
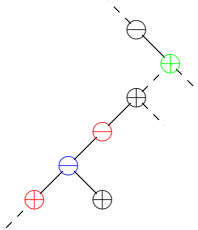
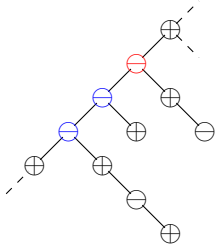
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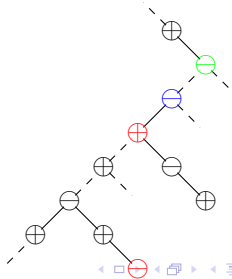
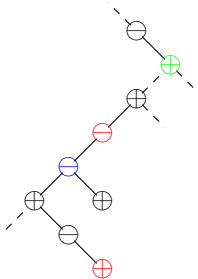
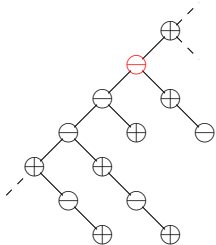
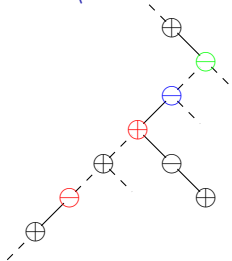
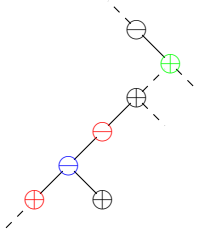
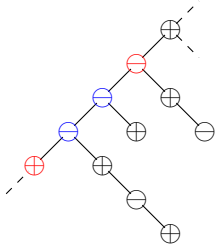
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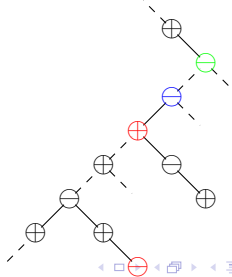
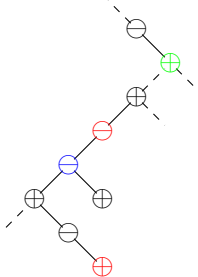
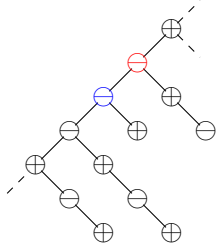
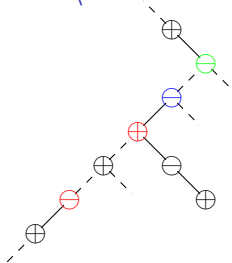
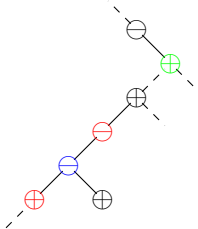
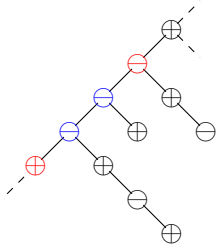
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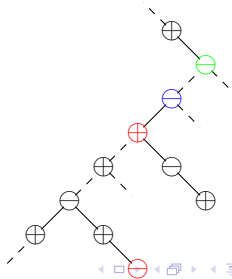
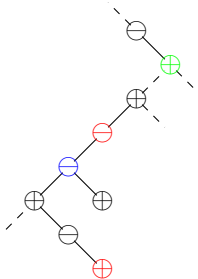
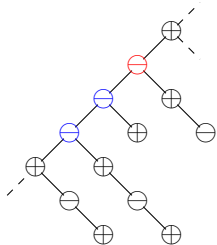
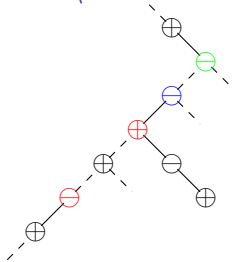
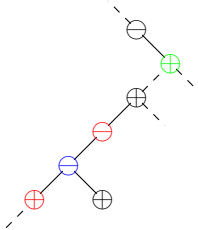
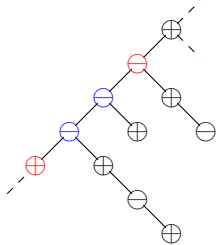
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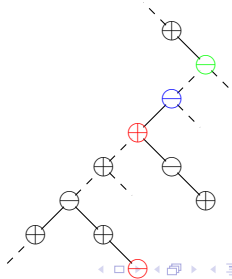
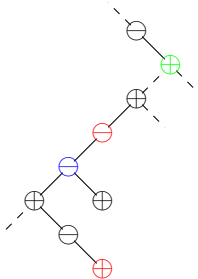
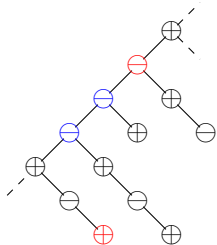
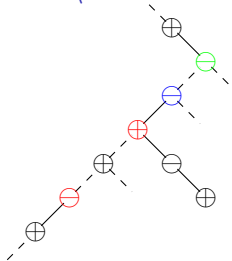
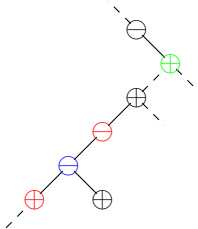
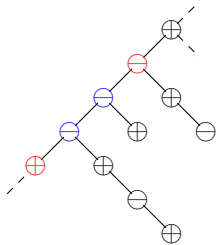
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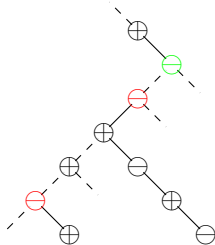
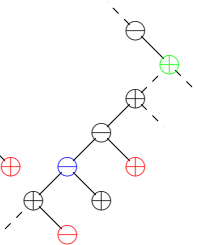
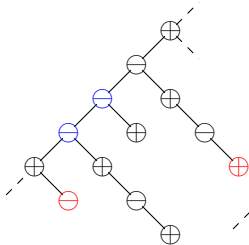
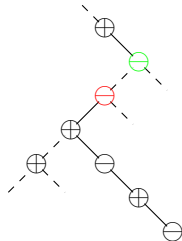
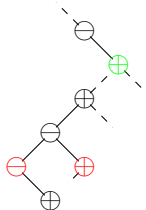
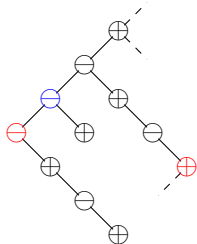
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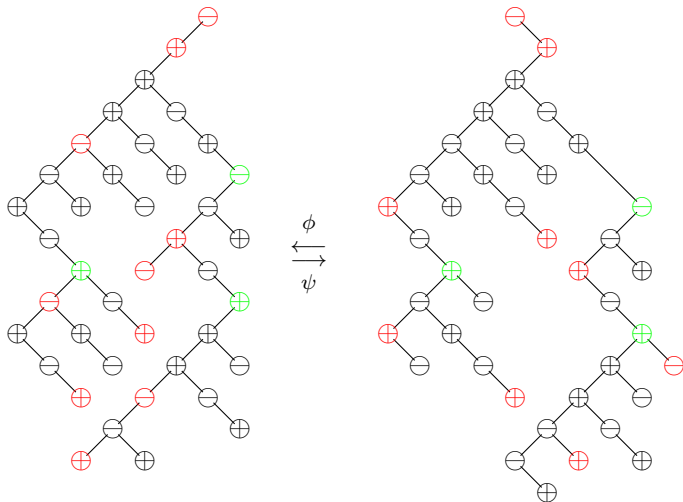
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Example with multiple cases



Outline

Introduction and Main Results

Hierarchy and Implication

Main Results

Separable permutation, Schröder word and di-sk tree

Direct sum and skew sum

Sweeping-algorithm

The di-sk tree

$S_n(t)$ is γ -nonnegative

Steps in the proof

The bijection

Conclusion

In retrospect. Both $\mathfrak{S}_n(312)$ and $\mathfrak{S}_n(3142, 2413)$ “respect” \mathfrak{S}_n 's γ -decomposition.

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$$S_n(t) := \sum_{\pi \in \mathfrak{S}_n(2413, 3142)} t^{\text{des}(\pi)} = \sum_{k \geq 0}^{\lfloor \frac{n-1}{2} \rfloor} \gamma_{n,k}^S t^k (1+t)^{n-1-2k},$$

$$\gamma_{n,k}^S = \#\{\pi \in \mathfrak{S}_n(3142, 2413) : \text{dd}(\pi) = 0, \text{des}(\pi) = k\}.$$

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The descent polynomials $S_n(t)$ and $D_n(t)$ are real-rooted for each $n \geq 2$.

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Thanks for listening!